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Turbo Macro Pro/TMPx Syntax

TMPx Assembler Overview

TMPx source code syntax is a superset of the x64-based sibling, Turbo Macro Pro. Basic elements of the syntax are:

labels: Something like variable names, a label is used to refer to a specific memory address or retain a value that can be used in calculations or assigned to a variable. Labels can be a memory address or just a numeric constant or multiply or any other expression that can be calculated.

numbers: These are the standard 6502 mnemonics, in the form of three letter abbreviations. TMPx observes 6502 mnemonics as well as illegal opcode mnemonics and recognizes specific built-in 65C02 or 65816 CPU (Central Processing Unit) and whose hardware.

pseudo-ops: Assembler directives that instruct TMPx to perform some operation instead of merely translating an opcode into machine code.

comments: A simple way to annotate source code; comments are completely ignored by the assembler.

Turbo Macro Pro -Assembler Details

Constant Values

Constant values can be expressed in either decimal, hexadecimal, binary or as characters depending on how the value is written. The largest constant value TMPx can recognize is $ffff (65,535).

```
;  denotes a hexadecimal value (limited to single bytes)
%  denotes a binary value (limited to single bytes)
'  denotes a character value, single character to a byte
```

Any value that is not preceded by $ or % or appended in quotes is treated as a decimal value.

```
=> Examples
$20          =  hex  $20, dec  32
$2000        =  hex $2000, dec 8192
%11111111    =  bin  %ff, dec  255
%1111111111  =  bin %3ff, dec 1023
'A'          =  hex  $41, dec  65
```

```
=> Examples
[petscii]    =  petscii code 91
[screen]     =  screen code 27
[127]        =  graphical character, PETSCII code 127
```

The loaded-number is further expanded to include any three digit decimal or hexidecimal value enclosed in curly braces. Such tokens are translated directly into the corresponding PETSCII code.

```
=> Examples
{$20}        =  space, PETSCII code 32
{%00000101}  =  PETSCII code 83
```

A loaded token can be used as a single character constant value enclosed in single quotes or as or or more characters in a double quoted string. Characters are translated into PETSCII codes. Control characters (such as a boot cat control) or characters (like decimal or hexidecimal). A full list of the named loaded tokens recognized by TMPx is given at the end of this documentation.

Labels

Valid labels in TMPx can be composed of any combination of letters, numbers, and the underscore character (obtained by pressing Shift-@). Labels must begin with a letter or underscore. Label names can be used in expressions (see below) in the same way as constant values.

To retain a label's value, it is assigned a constant value. Assign another value, or to an expression which will be resolved into an value. Once set, a label's value can be changed by assigning it a new value at the new address.

```
=> Examples
bah  = $2000    ;sets the label 'bah' to $1000.
lah = bah       ;sets the label 'lah' to equal whatever address bah
                ; represents, in this case $2000. If bah changes later,
                ; lah will change to match. Label 'lah' is equal the address that the
                ; expression 'bah+$' resolves to ($2000).
```

There is a special label, which is called the 'program counter' and is represented in source code as the asterisk (*). It is important to set the value of * at the top of your code. When TMPx assembles your source, * is used to tell the assembler where the machine code should be assembled in memory. Here are a few examples of how * is used:

```
=> Examples
*  = $1000   ;sets the program counter to $1000.
```

Note that unlike normal labels, you can reset the value of * anywhere you want in your source code. Each time you redefine *, the assembler will begin assembling machine code to the address specified by the new address.

Expressions

Expressions in TMPx can be formed with the following operators, all of which are binary operators (i.e. they require a value on the right AND left side of the operator):

```
+   for addition
-   for subtraction
*   for multiplication
/   for division
&   for bitwise and
|   for bitwise or
^   for bitwise xor
```

Values used with expressions can be constant values described above, or labels. Note that you can also use parentheses and as you would expect they affect the order in which the expression is evaluated.

```
=> Examples
$20 + 4          =  hex $24, dec  36
$1k - %00000011  =  hex $3fd, dec 1021
$2f&$f0          =  hex $20, dec  32
$f2|$f3          =  hex $f3, dec  51
$ff^$f0          =  hex $0f, dec  15
```

Unlike standard expressions in C/C++, there is no operator precedence. For example, multiplication has no higher precedence than addition.

```
[2+3]*4          =  hex $14, dec  20
2+3*4            =  hex $14, dec  20
```

In the first example above, the multiplication is resolved first, followed by the addition. In the next example, the addition is encountered first and is a numerical value, the multiplication follows. Although this is common practice in programming languages, it is important to keep in mind. Careful use of parentheses can help make the expected results explicit.

The special * can also be used in expressions, and it will always resolve to whatever the current program counter is where the assembler is outputting machine code to.

```
=> Examples
*  = $1000       ;sets the program counter to $1000.
jmp *+$05        ;assembles as jmp $1005
*  = $2000       ;sets the program counter to $2000 because the program counter
                 ; at this time will actually be $1005
```

The value of an expressions can also be modified by the following characters, which must appear at the very front of the expression:

```
<   denotes that the low byte of the following constant or
    expression should be taken
>   denotes that the high byte of the following constant or
    expression should be taken
!   denotes an expression whose value may currently or later
    be undefined. Normally, TMPx will flag this as an error
    but it should be remembered that it will be ignored in two
    of outputting machine code. See below for an example.
```

```
=> Examples
>$0114       =  hex $01
<$0114       =  hex $14
>$0114+$c20  =  hex $0d
```

Examples of the use of '!' are shown below:

```
lda $02     assembles as lda $02
lda !$02    assembles as lda $0002
```

Another use of the '!' is to prevent phase errors. A phase error occurs when the length of code or data changes and pass two of the assembler. Normally, TMPx will flag any label whose value has changed as a phase error.

```
    lda bah ; here, bah is undefined and TMPx assumes it is a word
bah = $12
bah = $02   ; bah is now defined as a byte value
```

Assembling the above code will result in a phase error, since TMPx assumes undefined labels represent words during the first pass. If the label is then found defined as a byte value, or if it actually resolves to a byte value, then the second pass will be shorter than the first and a phase error results. To prevent the phase error change the code to look like:

```
    lda !bah
bah = $12
bah = $02
```

The '!' tells TMPx that the expression following should be treated as a word no matter how the label was defined. Obviously this is a dangerous practice, but sometimes necessary. Remembering this behind-the-scenes behavior this is sometimes undesired so using '!' notation will help avoid the woes of phase errors.

Comments

You can add comments to your source code by using the ';' (semicolon). Any ';' found by TMPx will cause the assembler to ignore whatever follows the ';' til the end of that line. So comments can follow an opcode, or a pseudo-op, or they can be on a line by themselves.

```
=> Examples
;  this is a comment
lda #$d7   ; this is a comment
lda $2000  ; so is this!
```

Pseudo-ops: Data

The following pseudo-ops are the standard way to make data tables. Note that values are not necessarily constrained to constants - they can be complete expressions as above.

```
.byte   ;or .by  takes a byte string (which can be decimal, hex, binary
         or character constants and produces a table of bytes.
.word   ;takes a list of word values (separated by commas) and
         produces a table of bytes where the words are arranged in
         the standard 6502 low/high order.
.fcs    ;takes a list of word values and creates a table of bytes
         produces a table of bytes where the words are decremented
         by 1 except for the last word. This is used to create
         string which manipulation with a.r. return addresses.
```

```
=> Examples
.byte 25,'A',$ce  =  $19 $41 $ce
.word $1f2f       =  $2f $1f
.word $1f2f,$3f4f =  $2f $1f $4f $3f
```

A similar set of pseudo-ops can be used to make string tables. Note that loaded tokens can be used inside strings.

```
.text    takes a quoted string and converts it into a span of bytes.
         Characters are converted from ASCII to PETSCII including any
         loaded token conversion.
.null    ;sets a null terminated span of bytes is terminated with
         a null ($0) byte.
.shift   ;as .text, but the last byte has its high-bit set to 1.
.scr     ;takes the characters and converted into screencodes.
```

```
=> Examples
.text "hi"       =  $48 $49
.null "hi"       =  $48 $49 $00
.shift "hi"      =  $48 $c9
.scr "hi"        =  $08 $09
```

Another pseudo-op is provided for convenience when you want to specify a large span of repetitive values. .repeat with a count followed by a sequence of one or more values or expressions will create as many copies of the values/expression as you've specified.

```
.repeat ...produce 1 or more copies of a sequence of 1 or more values.
```

```
=> Examples
label = $3675

.repeat 8,$7f            =  $7f $7f $7f $7f $7f $7f $7f $7f
.repeat 3,label          =  $75 $36 $75 $36 $75 $36
.repeat 2,"A",$7f,$ff    =  $41 $42 $43 $45 $42 $43
```

Pseudo-ops: Assembler Control

These pseudo-ops allow TMPx to assemble chunks of code based on the evaluation of an expression - if 0 is equal or not equal to zero or if it is greater than zero. They can be used as the functionality to selectively assemble code based on the value of a label, for example, or even the current value of the program counter.

```
.if      tests value of expression for inequality with zero.
         (same as .ifne)
.ifne    tests value of expression for inequality with zero.
         (same as .if)
.ifeq    tests value of expression for equality with zero.
.ifpl    tests value of expression for plus (or zero) condition passes
.ifmi    tests value of expression for minus condition passes
         if the evaluated expression is from $8000-$ffff).
.ifnz    tests value of expression for non-zero condition passes
.ifze    tests value of expression for zero condition passes
         if the evaluated expression is from $0-$7fff).
.endif   marks the end of the code block that is assembled when the
         conditional test is met
```

```
=> Examples
cycle = 0

.ifeq cycle<$5 ;if cycle != $5
   nop         ;produce one nop
.endif

.ifeq *-$1000  ;if label has been defined,
   nop         ;produce one nop
.endif
```

In addition to conditions based on the value of an expression, two other pseudo-ops test based simply on whether a label had been defined or not, regardless of its value. Here are the pseudo-ops and examples of their use:

```
.ifdef   ;tests whether a label has been defined, condition passes if it has
.ifndef  ;tests whether a label has been defined, condition passes if it has not
```

```
=> Examples
*=$c000
.ifdef bah
   lda bah
.endif
```

Note that .ifdef and .ifndef can be used for testing whether a label definition has been specified in the command line arguments. As a scenario where this may be applied would be activating debugging code in output based on passing TMPx a "-DDEBUG" command line argument. For more information refer to the source with "RAM DEMO".

Pseudo-ops: Blocks (Scope Control)

Blocks in TMPx are used to provide useful separation: marking a section of code so that it can contain its own local labels, which can be redefined outside of the block or within a different block. Using blocks when coding reusable subroutines can save time from worrying about not reusing the same label name twice. Each block is inserted into another source; a color-switch example that the labels used inside the subroutine block will conflict with any labs already defined in the other code. Note that blocks can be nested within other blocks, each one creating a local scope for labels defined inside.

```
.block   starts a code block
.bend    ends a code block
```

```
=> Examples
*=$c000
.block
   lda .loop ;lda defined as $02
   sta .loop ;sta defined as $04
sub .block   ;our subroutine label name, 'sub'
   ldx .loop ;ldx defined as $03 but only applies in the block
   stx .loop
   jsr sub   ;assembles as: sub: $7f, not sub: $02!
   rts
.bend
   lda .loop ;assembles as: lda $02 again...
```

Pseudo-ops: Variables

Variables are a special kind of label that can be redefined (even without worrying about localized blocks). This is the only way to change the value of a label in the same context without getting a 'double defined' error. These are most useful when combined with other pseudo-ops like .repeat (described above) and .goto (see below).

```
.var   defines a variable identified by the label preceeding '.var'.
       which is set to the value of the expression following the '.var'.
```

```
=> Examples
va   .var $01   ;defines a var named va with value of 1
lda #va         ;results in lda #$01
va   .var va+1  ;redefines va to it's current value+1
lda #va         ;results in lda #$02
```

Pseudo-ops: Unconditional Goto

These pseudo-ops have the assembler to jump to a given label and continue assembling from there. This can be extremely useful for repeating a short code block 10, 100, or even more times. Only labels flagged with a .to pseudo-op can be reserved by a .goto!

```
lab   identifies a label that can be referenced by a .goto
.goto  causes assembler to goto the label referenced and continue
       assembling from that point on.
```

```
=> Examples
*  = $0100  ;set program counter to $0100
loop .to    ;label 'loop' is prepared to be a .goto target
     inx
     cpx #$ff ;.when x=$ff
     .free on1  ;if cnt != 0
     .goto loop ;then goto loop
     .endif     ;until x=$ff $100 nop's are assembled
```

Pseudo-ops: Macros

Macros are a means of inserting a large span of code by using a short name labelled (kind of) to be 6 arguments. This "macro call" is expanded to the full code when the assembler encounters it and the arguments given. A macro call is identified by a '#' followed by the macro label.

```
.macro   starts a macro definition that will implicitly treat the expanded
         code block as though it were enclosed by .block/.bend.
         (see .mblock below)
.segment  similar to .macro definition but without any implied block
.mend    ends a macro definition
```

Macro arguments are substituted as numbers in the macro code with either a '\' (ASCII backslash) to mark a substitution or an '@' symbol to mark a string substitution.

```
=> Examples
poke .macro   ;start macro def
     lda #\2  ;the \2 marks an argument substitution, in this case the
              ;second argument will replace the \2
     sta \1   ;another substitution
     .mend    ;end macro def
```

In the source code, a macro call for the above macro definition would look like:

```
#poke $d020,0
```

Upon assembly the macro call is expanded, and substitutions are made (in this case \1 is replaced by $d020 and \2 is replaced by 0, resulting in the following assembly being generated to replace the macro call:

```
lda #$00
sta $d020
```

Here is a more complete example showing recursion with macros:

```
table .macro   ;start macro def
      .word @1,@2,@3,@4,@5,@6,@7,@8,@9 ;plus 1
      #tab   ;a macro call within a macro!
      .mend  ;end macro def

tab   .segment ;start macro def
set   .var cnt:1 ;decrement cnt
      .if cnt  ;if cnt != 0
      .word *  ;then recursively macro call itself
      .endif   ;end of the if
      .word cnt ;store cnt here with current value of cnt
cnt   .var cnt-1 ;decrement cnt
      .mend    ;end macro def

      #table $0,:cnt1 ;the macro to create a byte table from 0 to 64
```

Here is an example of using a text parameter in a macro definition and call:

```
error  .macro    ;start macro def
       lda #<\1  ;lo byte of 1st arg
       ldy #>\1  ;hi byte of 1st arg
       jsr $ab1e ;print routine
       .byte @2  ;the text data
rts    .null "?\2 error" ;@2 is replaced by a textual argument
       .mend     ;end macro def

       #error "Bnc too far",bnc
```

So you see, use \x for numerical arguments, and @x for text arguments, where x denotes the argument number.

Test arguments can be even used to modify an opcode! Example:

```
do    .segment
      ldx #0
loop  .dnum
;
while .segment   ;start macro def
      bit loop   ;the bit will be replaced by a test argument
      .bend      ;end segment
      .mend      ;end macro def
;
      lda #5     ;start of code
      jsr while  ;call macro do
      dex
      lda #"*"
      jsr $ffd2  ;call macro do
      dex
      bne $f582  ;call macro while
      lda #"\n" ;call macro 'while' with parameter "\n"
      jsr $ffd2
      lda #"*"
      jsr while "\n" ;call macro 'while' with parameter "\n"
```

Pseudo-ops: Includes

The include pseudo-op allows you to assemble code from a different source file. For example, you could use .include to assemble a file with label definitions for all of the kernal jumptable, or you could write small reusable subroutines separately and include them in as you need them.

```
.include   ;load and assemble the specified file
```

```
=> Example
.include "kernal.s"
```

The .binclude pseudo-op will include the binary data in the given file directly into the assembled output at the location of the .binary call. You may optionally specify a number of data bytes to skip from the start of the included file by specifying that after the filename. This would be used to, for example, skip a two byte load address at the start of the file.

```
.binary   load and output bytes from the specified file, optionally
          skipping bytes at the start of the file
```

```
=> Example
.include "bm.dat"
.binary "music.prg",2
```

Pseudo-ops: Code Offsets

The .offs pseudo-op is used is often for memory address to where code is being assembled. This sounds identical to the program counter, but there is a subtle difference. The program counter (*) is where the assembler is outputting machine code. Typical code is a program counter of $1000 and the .offs is $c000. This means that the assembler is outputting machine code to $1000 but is assembling the code as though it were at $c000. For example, if you have a label defined as the address in the program counter but .offs changes the "actual address" to where object code is assembled. If you do not use calls to .offs in your code, then the actual address and effective address will always be identical. Some of the common reasons one may find it will remain properly in a different part of memory than where it will be executed are:

```
.offs   ;change the location that code is assembled to; code will be assembled forward
        ; to memory the the value of the expression following .offs
```

```
=> Example
(this code)  =  (assembles as)
*  = $1000
start  bit $2000     bit $2000
       lda #$00      lda #$00
       sta $2000     sta $2000
       rts           rts
.offs $c000
bank   lda #$20      lda #$20
       jmp start     jmp $2002
```

The code assembled from $2000 to $2002 by using an expression for the .offs as can generate code in a contiguous memory span, even though part of the code was assembled to execute at a different memory address (in this case, $0000). In this way, you can write code for different memory locations by assembling discontinued code locations in memory. Naturally this can be extremely useful for coding routines that need to execute in zero page, or for coding routines to run in drive memory.

Pseudo-ops: Printer Control

The printer pseudo-ops allow some control of output during assembly. .printf and .print allow you to skip printing sections of code, or even to change the verbosity of the control whether the expanded macro gets printed or just the macro call.

```
.prdo    turn on printing
.prdf    turn off printing
.hidemac  show unexpanded calls
.showmac  show expanded calls
```

Pseudo-ops: Miscellaneous

```
.scr     ;followed by name value, with which all code after the
         ;.scr call will be translated
.dsb     ;define storage block
.dsbnum  ;try = bounce 2,74,0,255,1,0,0,100,100,0,0 for a good time
```